Predicting Performance in an Introductory Programming Course by Logging and Analyzing Student Programming Behavior

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Abstract— The high failure rates of many programming courses means there is a need to identify struggling students as early as possible. Prior research has focused upon using a set of tests to assess the use of a student’s demographic, psychological and cognitive traits as predictors of performance. But these traits are static in nature, and therefore fail to encapsulate changes in a student’s learning progress over the duration of a course. In this paper we present a new approach for predicting a student’s performance in a programming course, based upon analyzing directly logged data, describing various aspects of their ordinary programming behavior. An evaluation using data logged from a sample of 45 programming students at our University, showed that our approach was an excellent early predictor of performance, explaining 42.49% of the variance in coursework marks – double the explanatory power when compared to the closest related technique in the literature.

Keywords- Learning Analytics, Prediction, CS1, Behavior.

I. INTRODUCTION

Due to a reputation for high failure rates [1], predicting a student’s performance in a first programming course is a well studied problem, and over the past fifty years, various predictors have been proposed. Early work mainly used standardized aptitude tests to predict performance [2]. As programming became more widespread, researchers (1980-90) began to explore a greater range of cognitive [3], psychological [4], and demographic [5] predictors. Researchers over the past two decades (1990-2010) extended prior work by exploring similar factors [6][7] and the predictive potential of new innovations in pedagogy [8][9].

However a limitation of studies to date is their tendency to use lengthy tests that often yield inconsistent results. Given potentially high enrollment numbers, the use of tests to gather predictive data can take a considerable amount of time for an instructor to process. Even if a test was indicative of performance, by the time it was processed, it may be too late for students to withdraw, or for instructors to intervene to prevent students from failing [7]. The criteria used for prediction is the main limitation of prior studies. Whilst cognitive, psychological, behavioral, or demographic traits may be indicative of performance, they are not directly related to the regular programming behavior of a student, or the programming tasks which they are required to perform. Because of these reasons, the indirect criteria used by prior studies fails to reflect changes in the learning progress and/or the learning behavior of a student over time.

There is a need to explore new predictors of performance, which are not based upon indirect criteria, but are instead based upon criteria which can be automatically measured, and directly reflect changes in a student’s learning progress. As well as being able to identify weaker students, such predictors could be used to drive an expert system [10][11] – providing weaker students with appropriate pedagogical interventions when required. A suitable measure could be based upon profiling a student by logging data describing various aspects of their ordinary programming behavior. Whilst recent research has provided visualizations of such data to instructors so that a manual intervention could be made [12], only [13] has attempted to collectively quantify several aspects of programming behavior into a predictor of performance. Jadud [13] proposes an algorithm called the Error Quotient (EQ) (revised in [14]). The algorithm uses a scoring function based upon the amount of errors a student encountered and how successive compilation failures in a session compare in terms of error message, location, and edit location. An overall score (range 0-1) for a student's performance during a session is computed by averaging the score of a set of successive compilation events. Higher EQ is indicative of weaker students. Although previously used by several studies [11][15][16], the EQ was shown to be a weak predictor of performance. This could be due to several methodological flaws concerning the incompleteness and inaccuracy of the approach, which we attempt to address and expand upon in our work (Sec. V). Our contributions include

- A unique approach for predicting performance based upon how a student responds to different types of error compared to their peers (proposing time as a predictor),
- A substantial improvement in terms of explanatory power and predictive accuracy by addressing the shortcomings of the main related approach [13][14].

II. ABOUT THE DATASET USED IN THIS STUDY

To explore possible predictors of achievement, we used a sample of students who studied the 2012/2013 Introduction to Programming (IP) course at our university. Programming behavior was directly logged by using an extension for the BlueJ IDE. Each time a student compiled their code on a university PC the extension would log a snapshot of their program source code along with the event type (success or fail), timestamp, error message reported and line number if applicable. Similar data was collected for invocations. As the use of final exams has been criticized as a means to
accurately measure programming ability [9], we use a student's overall coursework mark as the reference criterion of this study. This consisted of a weighting of their marks on a mid-term exam (25%), project (25%), practical exam (40%), and weekly lab exercises (10%). A total of 45 students (42 male) provided us with consent to use their logged data. 7 students indicated they had prior programming experience, but the majority indicated that the longest program they had written prior to course commencement was a medium length program (<2000 lines). Although data was logged over the duration of the course, due to the nature of student assignment work which involved intentionally propagating errors into source code, we restrict our analysis to the data gathered from 14 sessions (Term 1: weeks 3-9, Term 2: weeks 12-18).

III. The Watwin Algorithm
The uniqueness of our algorithm is to incorporate a scoring approach, where a student is relatively penalized based upon the amount of time that they take to resolve a specific type of error, compared to the resolve times of their peers. In the first stage of our algorithm, logged programming behavior is used to construct a set of successive compilation-event pairings, so that a student’s responses to different errors can be analyzed. This requires constructing consecutive pairings for each file that a student has attempted to compile, and estimating the amount of time a student has spent working on an error. In the second stage, each pairing is scored by assigning penalties based upon aspects of behavior which previous and our own research has identified as indicative of weaker performing students. Our algorithm is outlined as:

**Input:** A set of student programming logs (compilation and invocation) for all files a student compiled during a session.

1. **Prepare** a set of compilation pairings using the process presented in Sec. III (A).

2. **Quantify Programming Behavior**
   - **Score** each compilation pairing produced from (1) by using the scoring algorithm (Fig. 1).
   - **Normalize** each score by dividing by 35 (the maximum possible score for each pairing).
   - **Average** the normalized scores of all pairings.

**Output:** The mean average of all pairings (in the range 0-1), which is taken as the student’s Watwin score for the session. A score of 0 indicates that the student encounters no errors over a session. A score of 1 indicates that every compilation ended in an error, and that the student spent substantially longer than their peer’s between successive compilation events. The closer the score is to 0, the stronger the student.

A. Preparing a Set of Compilation Pairings

1) **Pair Construction.** For each file that a student attempted to compile during a session first construct a tuple of pairings \( \{e_1, e_2\}, \{e_2, e_3\}, \ldots, \{e_m, e_n\} \), using the compilation events associated with a file, ordered by timestamp. A naïve way to construct pairings would be to use the natural order that events occurred during a session. But this would fail to take into account the possibility of a student working on multiple files simultaneously, and can lead to an inaccurate representation of their programming behavior. For example, in a pairing \( \{e_i, e_j\} \) where \( e_i \) and \( e_j \) represent compilations of two distinct files, if the event type of \( e_i \) was ‘fail’ and the type of \( e_j \) was ‘success’, then the pairing \( \{e_i, e_j\} \) would incorrectly convey the student resolved the error of \( e_i \).

2) **Pair Pruning.** Identify and remove all pairings \( \{e_i, e_j\} \) where the code snapshots of \( e_i \) and \( e_j \) are identical. These cases can be caused by a ‘compile project’ feature of development software, and can artificially inflate the total number of compilation pairings. To take into account superficial changes which may have been made between compilations, such as adding comments or modifying layout, we first remove comments from the snapshots of \( e_i \) and \( e_j \) by using a regex expression. A standardized layout is then applied to the snapshots and compared for a match. If the snapshots are identical then \( \{e_i, e_j\} \) is removed. Also remove pairings where the event type of \( e_j \) was ‘success’.

3) **Filtering Commented and Deletion Fixes.** Whilst deleting and commenting code blocks can yield compilable files, these strategies provide little evidence of a student’s understanding of how to repair the actual fault. These actions can also be performed quickly; therefore the time taken to resolve an error in this manner may not be representative of the time taken to resolve using an actual fix. Deletion fixes are detected by computing the diff ratio between the snapshots of \( e_i \) and \( e_j \). If the count of insertions and changes = 0, and deletes > 0, then the pair is removed. Commented fixes are detected and removed by extracting the region of code surrounding the error location of \( e_i \), and using a regex expression to determine if the same fragment has only become commented in the snapshot of \( e_j \).

4) **Error Message Generalization.** Error messages within each compilation event pairing \( \{e_i, e_j\} \) are generalized by removing all identifier information. This allows us to build a profile for different classes of error, rather than for single specific messages. For example, “unknown class - Pet” becomes generalized to “unknown class”.

5) **Time Estimation.** The final step involves estimating the amount of time that a student has spent working on each compilation pairing \( \{e_i, e_j\} \). The simplest approach would be to directly compute the difference between timestamps of \( e_i \) and \( e_j \). But as our pairings are constructed on a per-file basis, this would fail to take into account whether a student has spent time working on other files between \( e_i \) and \( e_j \). We therefore first construct a combined sequence of invocation and compilation events \( \{h_1, h_2, \ldots, h_k\} \) for all files in a session, ordered by timestamp. For every \( \{e_i, e_j\} \), if there exists an \( h_k \) such that the timestamp of \( e_i > h_k > e_j \), we estimate the time spent on \( \{e_i, e_j\} \) as the difference between the timestamps of \( e_i \) and \( h_k \). The assumption is that a student has stopped working on the source of \( e_i \), and has instead only worked on the source code associated with \( h_k \).
B. Quantifying Ordinary Programming Behavior

1) Identifying Appropriate Predictors. Before developing a mechanism to meaningfully quantify a student's behavior, we first had to determine which aspects could indicate they were struggling to produce syntactically valid code. Prior research by [13] suggests that behavior exhibited by weaker students includes producing compilation pairings where both events result in compilation failures, have the same generalized message, and have the same error location. [15] found significant correlations between types of compilation pairings and performance where both events resulted in compilation failure and marginal correlations for pairings with the same message. Using our dataset we performed similar studies by correlating the average number of specific types of pairings which a student produced during a session, with performance. Significant correlations were found the average number of pairings whose both event types were compilation failures \((r(45)=-.43, p<.01)\), the average number of pairings where the generalized error message was the same \((r(45)=-.47, p<.01)\) and the average number of pairings where the generalized error messages were different \((r(45)=-.39, p<.01)\). Correlations were also found between average number of pairings with the same error location \((r(45)=-.26, p<.01)\). Our findings are consistent with [13] and [15], indicating that stronger programmers are associated with making less repeated errors, and will usually succeed in resolving an error in the next compilation.

A predictor which previous research [2-9][13-16] has not explored is the amount of time which a student takes to resolve an error. Research by [17] showed that the resolve times of certain types of errors vary based upon student ability; however, they did not use this variable to predict performance. We also hypothesized that in addition to having a higher frequency of errors; weaker students would also take longer to resolve errors than stronger students. After first removing outliers using the \(2MAD\) rule [18], we found a strong significant correlation between a student’s mean resolve time and performance \((r(45)=-.53, p<.01)\), which would seem to confirm our hypothesis. Because of this, we have incorporated resolve time as a predictor in our scoring model. However, different types of error can be more difficult for a student to resolve than others. Grouping the resolve times into 7 distinct classes of error (syntax, computation, identifiers, scope, exceptions, inheritance, abstraction) [11], a non-parametric Kruskal-Wallis test [18] confirmed that resolve times were significantly different between different classes of error \((\chi^2(6)=1512.88, p<.01)\).

2) Scoring Programming Behavior. Based upon these findings, instead of considering the amount of time that a student takes to resolve any error, we consider the time they take to resolve a generalized type (Sec. III A(4)) of error, in comparison to a distribution of resolve times of their peers. As these distributions are generally positively skewed, we use the robust \(2MAD\) approach [18] to remove outliers, and apply a penalty based upon where a student’s resolve time lies in the distribution. If their resolve time is more than one deviation below the mean, then they have resolved an error much faster than their peers - so we apply a low penalty. If a student’s resolve time is more than one deviation above the mean, then they have resolved an error much slower than their peers - apply a higher penalty. Otherwise, apply a mid-range penalty. The main advantage of scoring students in this manner is that we can implicitly take into account the relative difficulty of different types of error. For instance, suppose a student resolved a GUI error in 30 seconds. Compared to their peers, this may be a good time, and the student would incur a low penalty. However, if they took 30 seconds to resolve a ';' expected error, then compared to their peers, this may be a bad time, and the student would incur a higher penalty. After scoring all pairings using the scoring algorithm (Fig. 1), the scores of all pairings are normalized and averaged to produce a Watwin score.

3) Deriving Fair Penalties. The penalties assigned in the scoring algorithm (Fig. 1) were not determined through random guesswork. We first experimented by weighting the penalties of each component based upon the strengths of their correlations with performance. But, this produced a narrow range of Watwin scores, and we felt that a better spread of individuals was required. We therefore carried out a brute-force search of the space surrounding the parameters we had originally chosen. The regression models generated were ranked based upon their explanatory power, and penalties were then determined by repeated random sub-sampling of the strongest 100,000 results. Although not yielding the strongest possible explanatory model for our dataset, the derived parameters had the advantage of spreading the Watwin scores whilst simultaneously reducing the deviation between a student’s session scores. Along with the cross-validation we performed (Sec. IV), this supports the generalizability of our approach to independent datasets.

![Figure 1. Watwin Scoring Algorithm.](image-url)
IV. RESULTS AND EVALUATION

To evaluate the effectiveness of our algorithm as a predictor of a student’s programming performance, we performed a linear regression, using a student’s Watwin score as the independent variable, and their overall coursework mark as the dependent variable. We also considered the ability of Watwin as a classifier of student performance, based upon undergraduate degree boundaries set at our university (first \(\geq 70\%), second 50-69\%, third 40-49\%, and fail: <40\%).

An inspection of the scatter graph showed a linear relation existed between a student’s Watwin scores and performance, and that there were no significant outliers present. Residual independence was confirmed by the Durbin-Watson statistic (2.11), and the normality of residual distribution confirmed by an inspection of a histogram and P-P plot. We found that a linear regression based upon a student’s Watwin score could significantly predict performance, \(F(1, 43) = 31.77, p < .01\), explaining 42.49\% of the variance in coursework marks (a strong effect [19]). The final RMSE of the model was low at 6.91\% and the final accuracy of the predictive classifier was 75\%. Further validation of our model using leave-one-out cross validation yielded a mean \(R^2\) of .4204 (\(SD=0.13\)), RMSE of 7.09\% (\(SD=1.12\)), and classification accuracy of 75\% (\(SD=1.30\)), indicating a good level of consistency with the full model.

However, it is important to consider how our algorithm performs, in terms of accuracy and explanatory power over the duration of a course. Interestingly, previous work [2-9] [13-16], used all available data to drive their predictive models. But predicting a student’s failure at the end of a course leaves little time for an instructor intervention. Therefore for each session in both datasets, we computed a regression and the classification accuracy, using only the data which had been logged up to, or during the session.

We found that after 4 sessions, accuracy had risen into the 60’s range, and after 5 sessions accuracy leveled off and stayed in the 70’s range consistently over the duration of the course. However, measures of accuracy are reliant upon the underlying classification used. A more interesting analysis is to compare how the explanatory power of the regression changes over time. As can be seen from Fig. 2, by the end of the first term (week 9), a substantial percentage of the variance in coursework marks could be explained by our algorithm (30\%), which rose to over 40\% by the end of the second term. The average explanatory power of the algorithm was high, explaining 30.05\% (\(SD=15.97\)) of the variance in performance. This confirms that our approach is data driven, and performs less well when data is scarce.

| Data Sample Point | Watwin | | | Jadud | | | |
|-------------------|--------|---|---|--------|---|---|
|                   | \(R^2\) | RMSE | Acc. | \(R^2\) | RMSE | Acc. |
| End of Course     | .4249  | 6.91 | 75.56 | .1922  | 8.19 | 60.00 |
| Average           | .3005  | 7.60 | 68.83 | .1407  | 8.44 | 55.82 |

Figure 2. Explanatory Power of Watwin and Jadud During The Course

V. COMPARISON TO JADUD’S ERROR QUOTIENT

A. Addressing the Methodological Weaknesses

The major methodological flaw of Jadud’s Error Quotient [13] concerns the method used to construct a set of pairings. In Jadud’s work, a set of consecutive compilation pairings are created by using events in the order that they occurred during a session. As previously discussed, this approach is flawed, as it assumes that either students only work on a single source file, or work on multiple files in a linear manner. However, we have found that students do not work in this way, and switching between files is common. Using our dataset we built 45,001 compilation pairings using Jadud’s method. We found that 13,490 pairings (29.98\%) were based upon compilation events from two different files. This has serious implications for the validity of the approach. For instance, when examining pairings having event types in the form \{fail, success\}, we found that 2,138 (24.13\%) were based upon events from two different files. Almost 25\% of the cases indicated that a student had resolved an error, whereas in reality, they had simply compiled a different file. We addressed this shortcoming by constructing pairings on a per-filename basis, allowing us to more accurately profile student behavior based upon the evolution of code across distinct files.

Also, by constructing compilation pairings on a per-session basis, it is possible for the source code similarity to be calculated using the source of two distinct files meaning that extra compilation pairings will be included in the filtered set. There are no measures taken to check superficial changes made to source code can be incorrectly flagged as semantic changes. The flaws of the preparation and filtering methods have implications for the validity of the scoring algorithm used. In Jadud’s approach, pairings having event types in the form \{fail, success\} will score 0. But, it is possible that a large percentage of these pairings are invalid (30\% in our dataset). As a student’s error quotient is averaged using the sum of every pair from a session, having a large amount of invalid 0 scoring pairings can lower a student’s EQ, and inaccurately reflect their performance.

Finally, there are the fundamental differences between the Watwin and Jadud approaches to consider. Whilst we found that a student’s mean error resolve time strongly...
correlated with performance ($r(45) = -.53$), Jadud’s approach does not incorporate any scoring of behavior based upon this dimension. It also fails to take the type of error into account, and scores all errors equally. Very recent research [17] and this paper have both shown that students will find different types of error more difficult to resolve than others. Our uniqueness is to take these factors into account by relatively penalizing students based upon the amount of time they took to resolve an error, in comparison to a distribution of normal behavior defined by their peers.

### B. Evaluation of Performance

We applied Jadud’s algorithm to our datasets. Consistent with previous findings [13-17], we found Jadud’s EQ to be a weak predictor of performance, and that a student’s error quotient could explain less than half of the variance in performance, compared to their Watwin scores (Table 1). As can be seen from Fig 2 whilst the explanatory power of the EQ improves over time, it eventually levels off and remains a consistently weak predictor, only explaining between 15%-20% of the variance in performance over the final weeks of the course. This is also confirmed by the low standard deviations of average $R^2$ values of the EQ values (Table 1). In contrast, the explanatory power of the Watwin scores consistently increases over the duration of the course, and is a strong early predictor, explaining almost 30% of the variance in performance after 5-6 sessions of data has been collected. To explore the effect of the previously outlined methodological weaknesses of Jadud’s algorithm, we ran Jadud’s algorithm using pairings built using the Watwin algorithm. We found an increase in the explanatory power of Jadud’s model ($R^2 = .26 (+.07)$), suggesting that whilst an appropriate preparation technique can improve explanatory power, alone, it is not enough to match the performance of our scoring approach where students are relatively penalized based upon their resolve times and programming behavior.

### VI. Conclusion and Future Work

In this paper we presented Watwin, a dynamic algorithm designed to predict student performance in a programming course. Unlike prior work [2-9] which mainly used indirect criteria to predict performance, our approach is based upon analyzing directly logged, quantitative data describing aspects of a student’s ordinary programming behavior. This allows prediction of performance to evolve over time reflecting changes in the student’s learning progress without the need to use multiple tests that often yield inconsistent results. The originality of our algorithm is to incorporate a method, where a student is relatively penalized based upon the amount of time they took to resolve an error, in comparison to a distribution of normal behavior defined by the resolve times of their peers. We addressed the methodological weaknesses of the closest related approach [13-14], and an evaluation has shown that our approach is a good predictor of performance, even early in a course. Future work will aim to further validate our approach using data gathered from an independent sample of students, to identify more characteristics of programming behavior that are indicative of weaker students through the use of multivariate statistical [20] and data mining techniques [21], and to apply our algorithm within an expert system to select and supply appropriate compiler feedback to students [11].

### References